How to Save TERABYES of memory:

on-the-fly algorithms for 3- (and 4-) body forces in many-body systems

Collaborators:

W. Erich Ormand, Lawrence Livermore Plamen G. Krastev, Harvard Supercomputing Hai Ah Nam, SDSU/ Oak Ridge Collaborators-in-training: Joshua Staker & Micah Schuster, SDSU

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THE KEY ISSUES

Sparsity/storage requirements of matrices

Redundancy of matrix elements

Factorization and reduced storage of operations

Parallel distribution of operations

What's new in BIGSTICK

SPARSITY AND MATRIX STORAGE

A SPARSE MATRIX, BUT....

Despite sparsity, nonzero matrix elements can require TB of storage

•Typical dimensions and sparsity

Nuclide	valence space	valence Z	valence N	basis dim	sparsity (%)	
²⁰ Ne	"sd"	2	2	640	10	
^{25}Mg	"sd"	4	5	44,133	0.5	
⁴⁹ Cr	"pf"	4	5	6M	0.01	
⁵⁶ Fe	"pf"	6	10	500M	2x10 ⁻⁴	
$^{12}\mathrm{C}$	N _{max} =8	6	6	600M	4x10-4	2-body force
$^{12}\mathbf{C}$	N _{max} =8	6	6	600M	2x10 ⁻²	3-body force

A SPARSE MATRIX, BUT....

Despite sparsity, nonzero matrix elements can require TB of storage

Nuclide	Space	Basis dim	matrix store
⁵⁶ Fe	pf	501 M	4.2 Tb
⁷ Li	N _{max} =12	252 M	3.6 Tb
⁷ Li	N _{max} =14	1200 M	23 Tb
¹² C	N _{max} =6	32M	0.2 Tb
¹² C	N _{max} =8	590M	5 Tb
¹² C	N _{max} =10	7800M	111 Tb
¹⁶ O	N _{max} =6	26 M	0.14 Tb
¹⁶ O	N _{max} =8	990 M	9.7 Tb

A SPARSE MATRIX, BUT....

Despite sparsity, nonzero matrix elements can require TB of storage

Nuclide	Space	Basis dim	matrix store (2-body)	matrix store (3-body)
⁴ He	N _{max} =16	6 M	0.2 Gb	12 Tb
⁴ He	N _{max} =20	39 M	3 Tb	270 Tb
⁷ Li	N _{max} =10	43 M	0.4 Tb	176 Tb
¹² C	N _{max} =6	32M	0.2 Tb	6.2 Tb
¹² C	N _{max} =8	590M	5 Tb	200 Tb

REDUNDANCY OF MATRIX ELEMENTS

A Sparse Matrix, but....

• How the Hamiltonian is represented

"occupation representation"

$$|\alpha\rangle = \hat{a}_{n_1}^+ \hat{a}_{n_2}^+ \hat{a}_{n_3}^+ \dots \hat{a}_{n_N}^+ |0\rangle$$

n _i	1	2	3	4	5	6	7
α=1	1	0	0	1	1	0	1
α=2	1	0	1	0	0	1	1
α=3	0	1	1	1	0	1	0

$$\hat{H} = \sum_{ij} T_{ij} \hat{a}_i^{\dagger} \hat{a}_j + \frac{1}{4} \sum_{ijkl} V_{ijkl} \hat{a}_i^{\dagger} \hat{a}_j^{\dagger} \hat{a}_l \hat{a}_k$$

Usually, each *i* represents single-particle states with good *j,m*, parity

RECYCLED MATRIX ELEMENTS

Only a fraction of matrix elements are unique; **most are reused.** Reuse of matrix elements understood through *spectator* particles.



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of nonzero matrix elements vs. # unique matrix elements

Nuclide	valence space	valence Z	valence N	# nonzero	# unique
²⁸ Si	"sd"	6	6	$26 \ge 10^{6}$	3600
⁵² Fe	"pf"	6	6	$90 \ge 10^9$	21,500

Nuclide	ab initio space	basis dim	# nonzero m.e.s	# unique	avg redundancy
⁴ He	$N_{max} = 16$	6M	$2 \ge 10^{10}$	109	18
$^{12}\mathrm{C}$	N _{max} =8	600M	6 x 10 ¹¹	$5 \ge 10^{7}$	10,000

FACTORIZATION ALGORITHMS

Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*

We work in an *M*-scheme basis:

Because J^2 and J_z both commute with **H**, one does not need *all* basis states, but can use many-body basis restricted to the same *M*.

This is easy because M is an additive quantum number so it is possible for a single Slater determinant to be a state of good M.

(It's possible to work in a *J*-basis, e.g. OXBASH or NuShell, but each basis state is generally a complicated sum of Slater determinants).

Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*

Because the M values are discrete integers or half-integers (-3, -2, -1, 0, 1, 2, ... or -3/2, -1/2, +1/2, +3/2...) we can organize the basis states in discrete *sectors*

Example: 2 protons, 4 neutrons, total M = 0

$$M_{z}(\pi) = -4$$

 $M_{z}(\upsilon) = +4$
 $M_{z}(\pi) = -3$
 $M_{z}(\upsilon) = +3$

$$M_z(\pi) = -2$$
 $M_z(\upsilon) = +2$

Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*

In fact, we can see an example of factorization here because all proton Slater determinants in one M-sector *must* combine with all the conjugate neutron Slater determinants

Example: 2 protons, 4 neutrons, total M = 0

$$M_z(\pi) = -4: 2 \text{ SDs}$$
 $M_z(\upsilon) = +4: 24 \text{ SDs}$ 48 combined $M_z(\pi) = -3: 4 \text{ SDs}$ $M_z(\upsilon) = +3: 39 \text{ SDs}$ 156 combined $M_z(\pi) = -2: 9 \text{ SDs}$ $M_z(\upsilon) = +2: 60 \text{ SDs}$ 540 combined

Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*

In fact, we can see an example of factorization here because all proton Slater determinants in one M-sector *must* combine with all the conjugate neutron Slater determinants

M _z (π) = -4: 2 SDs	M _z (υ) = +4: 24 SDs	48 combined
$egin{array}{c} \pi_1 angle \ \pi_2 angle \ igstarrow \end{array}$	$egin{array}{c c} v_1 angle \ v_2 angle \ v_3 angle \ v_4 angle \end{array}$	$egin{aligned} \pi_1 angle &oldsymbol{ u}_1 angle \ \pi_2 angle &oldsymbol{ u}_1 angle \ \pi_1 angle &oldsymbol{ u}_2 angle \ \pi_2 angle &oldsymbol{ u}_2 angle \ dots\ \dots\ dots\ dots$
	$ {m v}_{24} angle$	$ \pi_1 angle u_{24} angle$
	TRIUMF – Feb 2012	$ \pi_2 angle\! u_{24} angle$

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Reuse can be exploited using exact factorization enforced through additive/multiplicative quantum numbers



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Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*

Factorization allows us to keep track of all basis states without writing out every one explicitly -- we only need to write down the proton/neutron components

The same trick can be applied to matrix-vector multiply



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Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*



There are potentially 48×48 matrix elements But for H_{pp} at most 4×24 are nonzero and we only have to look up 4 matrix elements

Advantage: **we can store 98 matrix elements as 4 matrix elements** and avoid 2000+ zero matrix elements.

Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*

 $M_{z}(\pi) = -4: 2 \text{ SDs} \qquad M_{z}(\upsilon) = +4: 24 \text{ SDs} \qquad 48 \text{ combined}$ $\begin{vmatrix} v_{1} \rangle \\ |v_{2} \rangle \\ |\pi_{2} \rangle \qquad H_{pp} = \begin{pmatrix} H_{11} & H_{12} \\ H_{21} & H_{22} \end{pmatrix} \qquad \begin{vmatrix} v_{1} \rangle \\ |v_{2} \rangle \\ |v_{3} \rangle \\ |v_{4} \rangle \\ \vdots \\ |v_{24} \rangle$

Advantage: **we can store 98 matrix elements as 4 matrix elements** and avoid 2000+ zero matrix elements.

Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*

M _z (π) = -4: 2 SDs	M _z (υ) =	= +4: 24 SDs 48 combined
$ \begin{vmatrix} \pi_1 \\ \pi_2 \end{vmatrix} \qquad H_{pp} = \begin{pmatrix} H_{11} & H_{12} \\ H_{21} & H_{22} \end{pmatrix} $	$egin{aligned} egin{aligned} egi$	$\begin{split} H_{pp} \pi_{1}\rangle \nu_{1}\rangle &= H_{11} \pi_{1}\rangle \nu_{1}\rangle + H_{12} \pi_{2}\rangle \nu_{1}\rangle \\ H_{pp} \pi_{2}\rangle \nu_{1}\rangle &= H_{12} \pi_{1}\rangle \nu_{1}\rangle + H_{22} \pi_{2}\rangle \nu_{1}\rangle \\ H_{pp} \pi_{1}\rangle \nu_{2}\rangle &= H_{11} \pi_{1}\rangle \nu_{2}\rangle + H_{12} \pi_{2}\rangle \nu_{2}\rangle \\ H_{pp} \pi_{2}\rangle \nu_{2}\rangle &= H_{12} \pi_{1}\rangle \nu_{2}\rangle + H_{22} \pi_{2}\rangle \nu_{2}\rangle \\ \vdots \\ H_{pp} \pi_{1}\rangle \nu_{24}\rangle &= H_{11} \pi_{1}\rangle \nu_{24}\rangle + H_{12} \pi_{2}\rangle \nu_{24}\rangle \\ H_{pp} \pi_{2}\rangle \nu_{24}\rangle &= H_{12} \pi_{1}\rangle \nu_{24}\rangle + H_{22} \pi_{2}\rangle \nu_{24}\rangle \end{split}$

Advantage: we can store 98 matrix elements as 4 matrix elements and avoid 2000+ zero matrix elements.

Reuse can be **exploited using exact factorization** enforced through *additive/multiplicative quantum numbers*

Comparison of nonzero matrix storage with factorization

Nuclide	Space	Basis dim	matrix store	factorization
⁷ Li	N _{max} =12	252 M	3600 Gb	96 Gb
⁷ Li	N _{max} =14	1200 M	23 Tb	624 Gb
¹² C	N _{max} =6	32M	196 Gb	3.3 Gb
¹² C	N _{max} =8	590M	5000 Gb	65 Gb
¹² C	N _{max} =10	7800M	111 Tb	1.4 Tb
¹⁶ 0	N _{max} =6	26 M	142 Gb	3.0 Gb
¹⁶ 0	N _{max} =8	990 M	9700 Gb	130 Gb

Comparison of nonzero matrix storage with factorization

⁴He

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{max} =14	2M	46 Gb	1.2 Gb	2 Tb	16 Gb
N _{max} =16	6M	200 Gb	4 Gb	12 Tb	60 Gb
N _{max} =18	16M	820 Gb	11 Gb	60 Tb	190 Gb
N _{max} =20	39M	3 Tb	29 Gb	270 Tb	600 Gb
N _{max} =22	86M	9 Tb	70 Gb	1.1 Pb	1.4 Tb

Comparison of nonzero matrix storage with factorization

⁴He

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{shell} =8	29 M	1.4 Tb	0.6 Gb	120 Tb	11 Gb
N _{shell} =9	93 M	8 Tb	1.7 Gb	870 Tb	40 Gb
N _{shell} =10	270 M	36 Tb	5 Gb	5 Pb	120 Gb
N _{shell} =11	700 M	150 Tb	12 Gb	28 Pb	350 Gb
N _{shell} =12	1.7 G	500 Tb	27 Gb	130 Pb	900 Gb
N _{shell} =13	4 G	1.7 Pb	60 Gb	500 Pb	2 Tb

Comparison of nonzero matrix storage with factorization

⁷Li

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{max} =8	6 M	36 Gb	1.5 Gb	1 Tb	26 Gb
N _{max} =10	43 M	430 Gb	10 Gb	170 Tb	250 Gb
N _{max} =12	250 M	4 Tb	60 Gb		

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{shell} =3	0.4 M	0.8 Gb	6 Mb	10 Gb	44 Mb
N _{shell} =4	45 M	330 Gb	0.3 Gb	9 Tb	4 Gb
N _{shell} =5	2 G	38 Tb	16 Gb	2 Pb	140 Gb
N _{shell} =6	50 G	2 Pb	87 Gb	170 Pb	3 Tb

Factorization

Comparison of nonzero matrix storage with factorization

⁹Be

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{max} =6	5 M	22 Gb	1 Gb	0.6 Tb	12 Gb
N _{max} =8	63 M	460 Gb	9 Gb	17 Tb	200 Gb
N _{max} =10	570 M	7 Tb	70 Gb		

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{shell} =3	4 M	15 Gb	30 Mb	240 Gb	240 Mb
N _{shell} =4	3 G	30 Tb	3 Gb	1 Pb	50 Gb
N _{shell} =5	400 G	12 Pb	130 Gb	800 Pb	3.6 Tb

Factorization

Comparison of nonzero matrix storage with factorization

 ^{10}B

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{max} =6	12 M	60 Gb	1.3 Gb	1.6 Tb	22 Gb
N _{max} =8	165 M	1.3 Tb	16 Gb	52 Tb	360 Gb

Comparison of nonzero matrix storage with factorization

¹²C

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{max} =6	32 M	170 Gb	3 Gb	5 Tb	60 Gb
N _{max} =8	590 M	5 Tb	45 Gb	200 Tb	1 Tb
N _{max} =10	8 G	100 Tb	440 Gb		

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{shell} =3	82 M	400 Gb	0.1 Gb	9 Tb	1.5 Gb
N _{shell} =4	600 G	10 Pb	43 Gb	580 Tb	0.9 Tb

Comparison of nonzero matrix storage with factorization

¹⁶O

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{max} =4	0.3 M	1 Gb	70 Mb	17 Gb	0.7 Gb
N _{max} =6	26 M	140 Gb	3 Gb	4 Tb	53 Gb
N _{max} =8	1 G	8.6 Tb	70 Gb		

Space	Basis dim	matrix store (2-body)	factorization (2-body)	matrix store (3-body)	factorization (3-body)
N _{shell} =3	800 M	6 Tb	0.7 Gb	140 Tb	7.5 Gb

Drawbacks of factorization/on-the-fly algorithms:

Much more complicated to code up (even matrix storage is not trivial)

Less flexible in basis—for example, importance truncation much harder (if even possible)

4-body is in principle straightforward

Experience in going from 2-body to 3-body shows most difficult part is correctly matching indices of input interaction to internal representation (+ induced phases etc) – useful to have *small* cases with known solutions for debugging

PARALLEL IMPLEMENTATION

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Factorization makes it easier to compute workload and distribute across multiple nodes



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PARALLEL IMPLEMENTATION

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THE BIGSTICK CODE

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Many-fermion code: 2nd generation after REDSTICK code (started in *Baton Rouge, La*.)

Arbitrary single-particle radial waveforms Allows local or nonlocal two-body interaction Applies to both nuclear and atomic cases

Runs on both desktop and parallel machines --can run at least dimension 100M+ on desktop (20 Lanczos iterations in 300 CPU minutes)

20-30k lines of codes Fortran 90 + MPI + OpenMP Partially funded by SciDAC Plans to run on 50,000-100,000 compute nodes Plans to publish code late 2012

THE BIGSTICK CODE

What's new since last year:

Full 3-body capability

Improved efficiencies in memory usage

OpenMP parallelization (WEO)

"2nd generation" MPI parallelization in progress

Looking for purveyors of 3-body interactions to partner with ! (Also, 4-body...)